

An efficient quantum algorithm for the hidden subgroup problem in extraspecial groups.

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The problem

- *Input*
 - A finite group G
 - A *hiding function* $f : G \rightarrow S$ which hides $H \leq G$
($f(x) = f(y) \Leftrightarrow xH = yH$).
- *Output*
 - *Generators for H .*

Theorem (Shor'94)

If G is Abelian then there is an efficient quantum algorithm using Fourier sampling and classical postprocessing:

- which finds H with probability $\geq 1 - 1/|G|$,
- in polynomial time in $\log|G|$.

A p -group is a finite group whose order is a power of p .

Definition

If G is a group,

- the **derived (commutator)** subgroup $G' = \langle \{[x, y] : x, y \in G\} \rangle$ and $[x, y] = x^{-1}y^{-1}xy$,
- the **center** $Z(G) = \{z \in G : gz = zg \text{ for all } g \in G\}$,

Property

An **extraspecial** group G is a p -group which satisfies:

- $G' = Z(G)$,
- $Z(G)$ is cyclic of prime order p ,

Application.

Quantum [error correction](#):

- Stabilizer codes: extraspecial 2-groups form the real part of the Pauli group,
- Clifford codes.

Efficiency.

If G is extraspecial then:

- $|G| = p^{2k+1}$
- an efficient algorithm should run in $\text{poly}(\log(|G|)) = \text{poly}(k \log(p))$, and both k and p can grow to infinity.

Efficient solutions for HSP:

- [Ivanyos,Magniez,Santha'01] when p is fixed constant,
- [Bacon,Childs,van Dam'05] Heisenberg group (case $k = 1$).

Extraspecial groups of size p^3 for odd p

There are **two extraspecial groups** of size p^3 :

- one of **exponent p** (all elements of the groups are of order p),
- one of **exponent p^2** (elements of the group are of order p or p^2).

H_p , Heisenberg group of exponent p .

$$\begin{aligned} H_p &= \left\{ \begin{pmatrix} 1 & a & b \\ 0 & 1 & c \\ 0 & 0 & 1 \end{pmatrix} ; (a, b, c) \in \mathbb{Z}_p^3 \right\} \\ &= \langle x, y, z ; x^p = y^p = z^p = 1, [x, y] = z, [x, z] = [y, z] = 1 \rangle \end{aligned}$$

A_p , application group of exponent p^2 .

$$\begin{aligned} A_p &= \left\{ \begin{array}{l} \mathbb{Z}_{p^2} \rightarrow \mathbb{Z}_{p^2} \\ t \mapsto at + b \end{array} \text{ where } a \equiv 1 \text{ modulo } p \text{ and } b \in \mathbb{Z}_{p^2} \right\} \\ &= \langle x, y, z ; x^{p^2} = y^p = 1, [x, y] = z = x^p, [y, z] = 1 \rangle \end{aligned}$$

$$G' = Z(G) = \{z^\ell \mid \ell \in \mathbb{Z}_p\}$$

Every element has a **unique representation** of the form

$$x^i y^j z^\ell, \text{ where } i, j, \ell \in \mathbb{Z}_p.$$

Extraspecial groups of size p^{2k+1} by generators

There are two extraspecial groups of size p^{2k+1} , one of exponent p , and one of exponent p^2 .

Extraspecial group of exponent p

$$H_p \times \dots \times H_p \text{ mod } z_1 = \dots = z_k,$$

It is generated by $2k + 1$ elements $x_1, y_1, \dots, x_k, y_k$ and z which satisfy the relations $x_i^p = y_i^p = z^p = 1$, $[x_i, y_i] = z$ for $i \in \mathbb{Z}_p$,

$$[x_i, x_j] = [x_i, y_j] = [y_i, y_j] = [x_i, z] = [y_i, z] = 1 \text{ for } i \neq j \in \mathbb{Z}_p.$$

$$G' = Z(G) = \{z^\ell \mid \ell \in \mathbb{Z}_p\}$$

Elements have a unique representation of the form:

$$x_1^{i_1} y_1^{i'_1} \dots x_k^{i_k} y_k^{i'_k} z^\ell, \text{ where } i_1, i'_1, \dots, i_k, i'_k, \ell \in \mathbb{Z}_p.$$

Example

For $k = 2$

$$x_1 y_1^2 y_2 \cdot x_1 y_1 x_2^2 z = x_1^2 y_1^2 y_2 z^2 \cdot y_1 x_2^2 z = x_1^2 y_2 z^2 \cdot x_2^2 z = x_1^2 x_2^2 y_2 z \cdot z = x_1^2 x_2^2 y_2 z^2$$

Extraspecial groups of size p^{2k+1} by generators

There are **two extraspecial groups** of size p^{2k+1} , one of **exponent p** , and one of **exponent p^2** .

Extraspecial group of exponent p^2

$$A_p \times H_p \times \dots \times H_p \text{ mod } z_1 = \dots = z_k,$$

It is generated by $2k$ elements $x_1, y_1, \dots, x_k, y_k$ which satisfy the relations $x_i^p = y_i^p = 1$, $[x_i, y_i] = x_1^p$ for $i \in \mathbb{Z}_p$, with $i \neq 1$, $x_1^{p^2} = y_1^p = 1$, and

$$[x_i, x_j] = [x_i, y_j] = [y_i, y_j] = 1 \text{ for } i \neq j \in \mathbb{Z}_p.$$

$$G' = Z(G) = \{z^\ell \mid \ell \in \mathbb{Z}_p\} \text{ where } z = x_1^p.$$

Elements have a **unique representation** of the form:

$$x_1^{i_1} y_1^{i_1'} \dots x_k^{i_k} y_k^{i_k'} z^\ell, \text{ where } i_1, i_1', \dots, i_k, i_k', \ell \in \mathbb{Z}_p.$$

Reduction lemma.

If f can be transformed into an efficient quantum hiding function for $HG' \leq G$ then the HSP has an efficient solution in G .

Hiding function lemma.

There is an efficient quantum hiding function for HG' .

- Groups of exponent p when p is large: **hard case**,
- Groups of constant exponent: **easy** (special case of previous),
- Groups of exponent p^2 when p is large: **reduces** to first case.

Theorem

We can solve the HSP in any extraspecial p -group.

Finding H is reducible to finding HG' .

Case 1: $G' \subseteq H$. Then $H = HG'$.

Case 2: $G' \cap H = \{1\}$.

- Then H is abelian since $h_1 h_2 = h_2 h_1 z^\ell \implies z^\ell \in G' \cap H$.
- Therefore HG' is abelian since G' is the center.
- The function f restricted to HG' hides H . This is an instance of the abelian HSP.



Example

$$G = H_3 = \{x^i y^j z^\ell : 0 \leq i, j, \ell \leq 2\}$$

$$H = \{1, x^2 y z, x y^2 z\}, \quad G' = \{1, z, z^2\}$$

$$HG' = \{1, z, z^2, x^2 y, x^2 y z, x^2 y z^2, x y^2, x y^2 z, x y^2 z^2\}$$

Finding HG' is reducible to a hiding function for HG' .

- For $g = x_1^{i_1} y_1^{j_1} \dots x_k^{i_k} y_k^{j_k} z^\ell$ let $\bar{g} = x_1^{i_1} y_1^{j_1} \dots x_k^{i_k} y_k^{j_k}$
- $\bar{G} = \{\bar{g} : g \in G\}$.
To make \bar{G} a group define $\bar{g}_1 * \bar{g}_2 = \overline{g_1 g_2}$.
- Then $(\bar{G}, *) \simeq (G/G', \cdot) \simeq \mathbb{Z}_p^{2k}$ is abelian and $HG' \cap \bar{G}$ is a subgroup of $(\bar{G}, *)$.
- The function hiding HG' in G hides also $HG' \cap \bar{G}$ in \bar{G} .
- $HG' = (HG' \cap \bar{G})G'$.



Example

$$G = \{x^i y^j z^\ell : 0 \leq i, j, \ell \leq 2\}$$

$$H = \{1, x^2 yz, xy^2 z\}, \quad G' = \{1, z, z^2\}$$

$$HG' = \{1, z, z^2, x^2 y, x^2 yz, x^2 yz^2, xy^2, xy^2 z, xy^2 z^2\}$$

$$\bar{G} = \{1, y, y^2, x, xy, xy^2, x^2, x^2 y, x^2 y^2\}$$

$$HG' \cap \bar{G} = \{1, x^2 y, xy^2\}$$

Coset state.

- For random $a \in G$

$$|aHG'\rangle = \frac{1}{\sqrt{p}} \sum_{i \in \mathbb{Z}_p} |aHz^i\rangle$$

- Group action of g on the coset state:

$$g \mapsto |aHG' \cdot g\rangle$$

- Eigenvalues:

$$\forall g = hz^t \in HG',$$

$$|aHG' \cdot g\rangle = |aHG'\rangle$$

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Perturbed coset state.

- For random $a \in G$ and $u \in \mathbb{Z}_p$

$$|aHG'_u\rangle = \frac{1}{\sqrt{p}} \sum_{i \in \mathbb{Z}_p} \omega^{-ui} |aHz^i\rangle$$

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$$\forall g = hz^t \in HG',$$

$$|aHG'_u \cdot g\rangle = \omega^{tu} |aHG'_u\rangle$$

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To **cancel** the disturbing phase: more sophisticated group action via group automorphisms.

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To **cancel** the disturbing phase: more sophisticated group action via group automorphisms.

Definition

For $j = 1, \dots, p-1$, automorphisms ϕ_j of G : $x_i \mapsto x_i^j$, $y_i \mapsto y_i^j$, $z \mapsto z^{j^2}$ for $i \in \{1, \dots, k\}$.

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- For random $a \in G$

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- Group action of g on the coset state:

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$$|aHG'_u\rangle = \frac{1}{\sqrt{p}} \sum_{i \in \mathbb{Z}_p} \omega^{-ui} |aHz^i\rangle$$

$$g \mapsto |aHG'_u \cdot \phi_j(g)\rangle$$

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- For random $a \in G$ and $u \in \mathbb{Z}_p$

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$$g \mapsto |aHG'_u \cdot \phi_j(g)\rangle$$

$$\forall g = hz^t \in HG',$$

$$|aHG'_u \cdot g\rangle = \omega^{u(j-j^2)t + uj^2t} |aHG'_u\rangle$$

To **cancel** the disturbing phase: more sophisticated group action via group automorphisms.

Definition

For $j = 1, \dots, p-1$, automorphisms ϕ_j of G : $x_i \mapsto x_i^j$, $y_i \mapsto y_i^j$, $z \mapsto z^{j^2}$ for $i \in \{1, \dots, k\}$.

Four copies state.

For $\bar{a} \in G^4$, $\bar{u} \in \mathbb{Z}_p^4$, $\bar{j} \in (\mathbb{Z}_p^*)^4$ and $g \in G$, let

$$|\Psi_g^{\bar{a}, \bar{u}, \bar{j}}\rangle = \bigotimes_{i=1}^4 |a_i HG'_{u_i} \cdot \phi_{j_i}(g)\rangle = \omega^{\sum_{i=1}^4 u_i((j_i - j_i^2)\ell + j_i^2 t)} \bigotimes_{i=1}^4 |a_i HG'_{u_i}\rangle.$$

Fact

$g \mapsto |\Psi_g^{\bar{a}, \bar{u}, \bar{j}}\rangle$ is *hiding* HG' for $\bar{j} \in (\mathbb{Z}_p^*)^4$ *solution* of

$$\begin{cases} \sum_{i=1}^4 u_i(j_i - j_i^2) & = 0 \\ \sum_{i=1}^4 u_i j_i^2 & = 0. \end{cases}$$

Lemma

$$\Pr_{\bar{u} \in \mathbb{Z}_p^4} [\exists \bar{j} \in (\mathbb{Z}_p^*)^4 \text{ solution}] \approx \frac{1}{2}.$$

The system is efficiently *solvable*.

Algorithm

- 1 Compute, for some $\bar{a} \in G^4$, the superposition

$$\frac{1}{p^2} \bigotimes_{i=1}^4 \sum_{u_i \in \mathbb{Z}_p} |u_i\rangle |a_i HG'_{u_i}\rangle$$

- 2 Measure the registers for \bar{u}
- 3 Solve the equation

$$\begin{cases} \sum_{i=1}^4 u_i(j_i - j_i^2) = 0 \\ \sum_{i=1}^4 u_i j_i^2 = 0 \end{cases}$$

- 4
 - If the system has a solution $\bar{j} \in (\mathbb{Z}_p^*)^4$ then

$$|g\rangle \mapsto \bigotimes_{i=1}^4 |a_i HG'_{u_i} \cdot \phi_{j_i}(g)\rangle$$

- else repeat 1., 2. and 3.

A **constant** number of repetitions is enough.

- We can generalise this method to nilpotent p -groups of class 2.
- Extend to higher classes.